

A characterization of the Taylor expansion of λ -terms

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Recall on quantitative semantics

λ -calculus

$$X \mapsto (F)X$$

$$\downarrow \beta^*$$
$$Y$$

Semantics

$$x \mapsto f(x)$$



Recall on quantitative semantics

λ -calculus

$$\begin{array}{c} X \mapsto (F)X \\ \downarrow \beta^* \\ Y \end{array}$$

Semantics

$$\begin{array}{c} x \mapsto f(x) \\ \parallel \\ \sum_{n=0}^{\infty} \frac{1}{n!} (\partial_x^n f \cdot x^n) 0 \end{array}$$

$\sum_{n=0}^{\infty} \frac{1}{n!} (\partial_x^n f \cdot x^n) 0$ is the *Taylor Expansion* of f

λ -calculus $\xrightarrow{\text{Taylor Expansion}}$ Resource calculus

λ -calculus

Grammar: $\Lambda : T, U ::= x \mid \lambda x.T \mid (T)U$

$$(\lambda x.T)U \xrightarrow{\beta} T[U/x]$$

Resource calculus

Grammar: $\Delta : t, u ::= x \mid \lambda x.t \mid \langle t \rangle [u_1, \dots, u_n]$

$$\langle \lambda x.t \rangle [u_1, \dots, u_n] \xrightarrow{r} \sum_{\sigma \in \mathcal{S}_n} t\{u_{\sigma(1)}/x_1, \dots, u_{\sigma(n)}/x_n\}$$

Substitutes each occurrence of x in t by only one u_i

Reduces to 0 otherwise

$$M \xrightarrow{\text{Taylor Expansion}} \sum_{t \in \text{taylor}(M)} \alpha_t t \xrightarrow{\text{NF}} \text{NF} \left(\sum_{t \in \text{taylor}(M)} \alpha_t t \right)$$

Goal: Characterize the image of this transformation



Theorem [Characterization]:

$\exists M$ λ -term s.t. $\sum_{t \in \Delta} \alpha_t \cdot t = \text{NF}(\text{taylor}(M))$ iff

0 ...

1 ...

2 ...

3 ...



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0 Theorem [Ehrhard - Regnier]:

$\forall \alpha_t \in \text{NF}(\text{taylor}(M)),$ if $\alpha_t \neq 0$ then $\alpha_t = \frac{1}{m(t)}$

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$\forall \mathcal{T} \subseteq \Delta, \exists M \lambda\text{-term s.t. } \mathcal{T} = \text{NF}(\tau(M)) \text{ iff}$

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1 $\text{FV}(\mathcal{T})$ is finite

2 \mathcal{T} is r.e.

3 ...

Conditions 1 and 2: based on Barendregt's theorem

Theorem [Barendregt]:

Let \mathcal{B} be a Böhm-like tree. There is a λ -term M such that $\text{BT}(M) = \mathcal{B}$ if, and only if, $\text{FV}(\mathcal{B})$ is finite and \mathcal{B} is r.e..



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Resource calculus and Taylor expansion

Ideal

Two corollaries and further works



Plan

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Two corollaries and further works

Resource calculus

Grammar: $\Delta : t, u ::= x \mid \lambda x.t \mid \langle t \rangle [u_1, \dots, u_n]$

Relation \xrightarrow{r} (strongly normalizing, confluent):

$$\langle \lambda x.t \rangle [s_1, \dots, s_n] \xrightarrow{r} \begin{cases} \{t\{s_{\sigma(1)}/x_1, \dots, s_{\sigma(n)}/x_n\} \mid \sigma \in \mathcal{S}_n\} \\ \emptyset \text{ if } \deg_x(t) \neq n \end{cases}$$

Unique normal form: $\text{NF}(t)$

$$\text{NF}(\mathcal{T}) \triangleq \bigcup_{t \in \mathcal{T}} \text{NF}(t)$$

Taylor expansion: $\Lambda \longrightarrow \mathcal{P}(\Delta)$

$$\tau(x) \triangleq \{x\}$$

$$\tau(\lambda x.T) \triangleq \{\lambda x.t \mid t \in \tau(T)\}$$

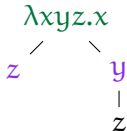
$$\tau((T)U) \triangleq \{\langle t \rangle [u_1, \dots, u_k] \mid t \in \tau(T); k \in \mathbb{N}; u_1, \dots, u_k \in \tau(U)\}$$



A first example : S

$$S := \lambda xyz.((x)z)(y)z$$

Böhm tree of S :



Taylor expansion of S:

$$\begin{aligned} \tau(S) &= \{\lambda xyz.\langle x \rangle 11, \lambda xyz.\langle x \rangle \underbrace{[z, \dots, z]}_n \underbrace{[\langle y \rangle 1, \dots, \langle y \rangle 1]}_m, \dots\} \\ &= \{\lambda xyz.\langle x \rangle [z^n][\langle y \rangle [z^{n_1}], \dots, \langle y \rangle [z^{n_k}]] ; k, n, n_1, \dots, n_k \in \mathbb{N}\} \\ &= \text{NF}(\tau(S)) \end{aligned}$$



Two other examples

$$(\mathbf{S})\mathbf{II} = ((\lambda x y z. ((x)z)(y)z)\lambda x. x)\lambda x. x \xrightarrow{\beta^*} \lambda x. (x)x = \delta$$

$$\tau((\mathbf{S})\mathbf{II}) = \{\langle \lambda x y z. \langle x \rangle 11 \rangle 11,$$

$$\langle \lambda x y z. \langle x \rangle [z, \dots, z][\langle y \rangle 1, \dots, \langle y \rangle 1] \rangle [1, \dots, 1][1, \dots, 1], \dots \}$$

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$$\mathbf{\Omega} = (\delta)\delta$$

$$\tau(\mathbf{\Omega}) = \{\langle \lambda x.\langle x \rangle [x^{n_0}] \rangle [\lambda x.\langle x \rangle [x^{n_1}], \dots, \lambda x.\langle x \rangle [x^{n_k}]]; k, n_0, \dots, n_k \in \mathbb{N}\}$$

$$\mathbf{NF}(\tau(\mathbf{\Omega})) = \emptyset$$



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Ideal

Terms in normal form: $\Delta^{\text{NF}} : t ::= \lambda x_0 \dots x_{m-1} . \langle y \rangle \mu_0 \dots \mu_{n-1}$

μ_i : finite multisets of simple terms in normal form

Uniform approximation \preceq

$\lambda x_0 \dots x_{m-1} . \langle y \rangle \mu_0 \dots \mu_{n-1} \preceq t$ iff

- (i) $t = \lambda x_0 \dots x_{m-1} . \langle y \rangle \nu_0 \dots \nu_{n-1}$
- (ii) $\forall i < n, |\mu_i| \neq \emptyset \implies \exists v \in |\nu_i|, \forall u \in |\mu_i|, u \preceq v$

\preceq -ideal

$\mathcal{T} \in \mathcal{P}(\Delta^{\text{NF}})$ ideal: downward closed, directed

- $\tau(\mathbf{S}) = \{\lambda x y z . \langle x \rangle [z^n][\langle y \rangle [z^{n_1}], \dots, \langle y \rangle [z^{n_k}]] ; k, n, n_1, \dots, n_k \in \mathbb{N}\}$
- $\{\langle x \rangle [y, z]\}$
- $\{x[x]\}$,



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- $\{\langle x \rangle [y, z]\}$
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Two corollaries and further works

Corollary 1

Let $\mathcal{T} \in \mathcal{P}(\Delta^{\text{NF}})$.

There is a normalizable λ -term M such that $\text{NF}(\tau(M)) = \mathcal{T}$ iff

- (i) $\text{height}(\mathcal{T})$ is finite
- (ii) \mathcal{T} is a maximal clique

$\Delta^{\text{NF}} : t ::= \lambda x_0 \dots x_{m-1} \langle y \rangle \mu_0 \dots \mu_{n-1}$, μ_i finite multisets of simple terms in normal form.

Coherence \supset on Δ^{NF} :

$\lambda x_0 \dots x_{m-1} \langle y \rangle \mu_0 \dots \mu_{n-1} \supset t$ iff

- (i) $t = \lambda x_0 \dots x_{m-1} \langle y \rangle \nu_0 \dots \nu_{n-1}$
- (ii) $\forall i < n, \forall u, u' \in |\mu_i \cdot \nu_i|, u \supset u'$

Clique: subset of a \preceq -ideal

$\mathcal{T} \in \mathcal{P}(\Delta^{\text{NF}})$ clique: $\forall t, t' \in \mathcal{T}, t \supset t'$

Corollary 2

Let $\mathcal{T} \in \mathcal{P}(\Delta^{\text{NF}})$.

There is a total λ -term M such that $\text{NF}(\tau(M)) = \mathcal{T}$ iff

1. $\text{FV}(\mathcal{T})$ is finite
2. \mathcal{T} is r.e.
3. \mathcal{T} is a maximal clique

Total terms

- (i) $M \xrightarrow{h^*} \lambda x_0 \dots x_{m-1}.(y)M_0 \dots M_{n-1}$
- (ii) M_0, \dots, M_{n-1} are total



Further works

Bring the results to more expressive calculi:

- $\Lambda\mu$ -calculus
 - Cannot use Barendregt's theorem
- Non-Deterministic settings
 - Cannot use Ehrhard Regnier's theorem
- ...